



RIEMANNIAN GAME DYNAMICS

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Outline

Background

Geometric Preliminaries

The Dynamics

Analysis



Game dynamics

Rationality concepts can often be justified as a the long-term outcome of simple dynamic processes (*game dynamics*). However, your mileage may vary:

- ▶ Replicator dynamics (Taylor and Jonker, 1978)
 - ▶ Positive correlation between payoffs and dynamics
 - ▶ Nash equilibria are fixed points – but so is every pure state
 - ▶ Dominated strategies are eliminated
 - ▶ Periodic orbits in 2-player zero-sum games
- ▶ Projection dynamics (Nagurney and Zhang, 1997)
 - ▶ Positive correlation between payoffs and dynamics
 - ▶ Nash \Leftrightarrow fixed points
 - ▶ Dominated strategies may survive
 - ▶ Limit cycles and periodic orbits in 2-player zero-sum games



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 - ▶ Limit cycles and periodic orbits in 2-player zero-sum games ✓✗

What this talk is about:

Provide a unifying framework for the foregoing dynamics and extend a range of rationality properties known for certain special cases.



Population games

Throughout this talk we will focus on *population games*:

- ▶ Continuum of *nonatomic players*
- ▶ Common (finite) set of *actions* \mathcal{A}
- ▶ *Population states* $x = (x_\alpha)_{\alpha \in \mathcal{A}}$ are defined by the proportions $x_\alpha \in [0, 1]$ of players playing each action $\alpha \in \mathcal{A}$.
- ▶ Set of population states: $\mathcal{X} = \Delta(\mathcal{A})$.
- ▶ Individual action rewards determined by associated *payoff functions* $v_\alpha: \mathcal{X} \rightarrow \mathbb{R}$

Example (Random matching)

- ▶ A population of agents is randomly matched to play a 2-player symmetric normal form game with payoff matrix $A = (A_{\alpha\beta})_{\alpha, \beta=1}^n$.
- ▶ The payoff to α -strategists when the population state is $x \in \mathcal{X}$ is

$$v_\alpha(x) = \sum_{\beta \in \mathcal{A}} A_{\alpha\beta} x_\beta$$



Evolutionary dynamics

Evolutionary dynamics are rules that assign to games a dynamical system of the form:

$$\dot{x} = V(x) \quad (\text{ED})$$

where $V(x)$ is usually specified as a function of state/payoff pairs: $V(x) \equiv \tilde{V}(x, v(x))$.



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Evolutionary dynamics are *admissible* if and only if $x(t) \in \mathcal{X}$ for all time $\implies V(x)$ must lie in the *tangent cone* of \mathcal{X} at x :

$$V(x) \in \text{TC}_{\mathcal{X}}(x) = \left\{ z \in \mathbb{R}^A : \sum_{\alpha} z_{\alpha} = 0 \text{ and } z_{\alpha} \geq 0 \text{ if } x_{\alpha} = 0 \right\}$$



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In many cases, $V(x)$ actually lies in the *tangent space* to \mathcal{X} at x , viz.:

$$V(x) \in \text{T}_{\mathcal{X}}(x) \equiv \left\{ z \in \mathbb{R}^A : \sum_{\alpha} z_{\alpha} = 0 \right\}$$

In that case, **no new strategies can enter the fray** (the support of $x(t)$ remains invariant)



Revision protocols

Microfoundations for (ED) are often provided by means of a *revision protocol* that specifies the (unconditional) *switch rates* $s_{\alpha\beta}$ at which α -strategists revise their strategy to β .



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Such protocols $s_{\alpha\beta}$ are usually written as

$$s_{\alpha\beta} = p_{\alpha} r_{\alpha\beta}$$

where

- ▶ p_{α} is the rate at which α -strategists receive a revision opportunity.
- ▶ $r_{\alpha\beta}$ is the (*conditional*) switch rate from α to β (i.e. once an α -strategist has received a revision opportunity).



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The *mean dynamics* generated by a revision protocol are then:

$$\dot{x}_{\alpha} = \sum_{\beta} [s_{\beta\alpha}(x, v(x)) - s_{\alpha\beta}(x, v(x))], \quad (\text{MD})$$

i.e the rate of change of x_{α} is the *inflow/outflow* difference from/to strategy α .



Exampel 1: the replicator dynamics

The *replicator dynamics* (Taylor and Jonker, 1978) are defined as

$$\dot{x}_\alpha = x_\alpha \left[v_\alpha(x) - \sum_\beta x_\beta v_\beta(x) \right] \quad (\text{RD})$$

i.e. the *per capita* growth rate of a strategy is the difference between its individual payoff and the mean population payoff.

Revision protocols that give rise to (RD) (up to constant payoff shifts):

- ▶ Imitation of success: $s_{\alpha\beta}(x, v) = x_\alpha x_\beta v_\beta$
- ▶ Pure imitation driven by dissatisfaction: $s_{\alpha\beta}(x, v) = -x_\alpha x_\beta v_\alpha$
- ▶ Pairwise proportional imitation: $s_{\alpha\beta}(x, v) = x_\alpha x_\beta [v_\beta - v_\alpha]_+$



Example 2: projection dynamics

The *projection dynamics* (Nagurney and Zhang, 1978; Lahkar and Sandholm, 2008) are defined on the interior of \mathcal{X} as:

$$\dot{x}_\alpha = v_\alpha(x) - \frac{1}{|\mathcal{A}|} \sum_{\beta \in \mathcal{A}} v_\beta(x) \quad (\text{RD})$$

i.e. the (absolute) growth rate of a strategy is given by the difference between its individual payoff and the (unweighted) average strategy payoff.

Revision protocols that give rise to (RD) (up to constant payoff shifts):

- ▶ Imitation of success driven by insecurity: $s_{\alpha\beta}(x, v) = \frac{1}{|\mathcal{A}|} v_\beta$
- ▶ Pure imitation driven by dissatisfaction and insecurity: $s_{\alpha\beta}(x, v) = -\frac{1}{|\mathcal{A}|} v_\alpha$
- ▶ Pairwise proportional imitation driven by insecurity: $s_{\alpha\beta}(x, v) = \frac{1}{|\mathcal{A}|} [v_\beta - v_\alpha]_+$

Similar to the replicator dynamics, but agents now quicker to abandon strategies that are used by few others.



Geometric underpinnings

- ▶ Under the projection dynamics, \dot{x} is the closest point projection of $v(x)$ to the tangent cone of \mathcal{X} at x – but what does “closest” mean?
- ▶ In potential games, the replicator dynamics ascend the game’s potential function, and this ascent is the steepest possible under a **non-Euclidean** metric
- ▶ In potential games, payoffs are potential gradients, so the replicator dynamics also represent a “closest point projection” – but under a **non-Euclidean** metric



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Objective of the talk

Study evolutionary dynamics whose vectors of motion $V(x)$ are aligned to payoffs $v(x)$ to the “greatest possible” extent.



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Scalar products and metrics

The notion of “closeness” presupposes a notion of “distance”. Given a vector space W , this is often defined via a *scalar product*, i.e. a bilinear pairing $\langle \cdot, \cdot \rangle$ that is:

1. **Symmetric:** $\langle w, w' \rangle = \langle w', w \rangle$ for all $w, w' \in W$
2. **Positive-definite:** $\langle w, w \rangle \geq 0$ for all $w \in W$, with equality iff $w = 0$.



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2. *Positive-definite*: $\langle w, w \rangle \geq 0$ for all $w \in W$, with equality iff $w = 0$.

The *norm* of $w \in W$ is then defined as

$$\|w\| = \langle w, w \rangle^{1/2}$$

or, in components w.r.t. a basis $\{e_\alpha\}_{\alpha=1}^n$ of W :

$$\|w\| = \left(\sum_\alpha \sum_\beta g_{\alpha\beta} w_\alpha w_\beta \right)^{1/2}$$

where the so-called *metric elements* $g_{\alpha\beta}$ of $\langle \cdot, \cdot \rangle$ are defined as:

$$g_{\alpha\beta} = \langle e_\alpha, e_\beta \rangle$$

Notation: $\langle \cdot, \cdot \rangle$ will often be identified with its *metric tensor* $g = (g_{\alpha\beta})_{\alpha,\beta=1}^n$



Riemannian metrics

Scalar products can be position-dependent, leading to the notion of a *Riemannian metric* :

Definition

A *Riemannian metric* on an open subset U of W is a continuous assignment of a metric tensor $g(x)$ to each point $x \in U$. In words, a Riemannian metric on U prescribes a way of measuring angles and lengths at each point $x \in U$, so it determines the *geometry* of U .



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Examples

1. The *Euclidean metric* on $U = \mathbb{R}^A$ is

$$g_{\alpha\beta}(x) = \delta_{\alpha\beta}$$

leading to the familiar expressions $\langle w, w' \rangle_x = \sum_{\alpha} w_{\alpha} w'_{\alpha}$ and $\|w\|_x = (\sum_{\alpha} w_{\alpha}^2)^{1/2}$

2. The *Shahshahani metric* on $U = \mathbb{R}_{>0}^A$ is given by

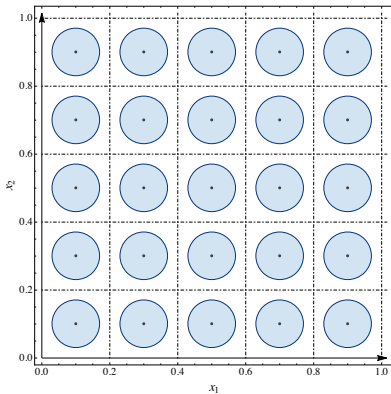
$$g_{\alpha\beta}(x) = \delta_{\alpha\beta}/x_{\beta}$$

leading to $\langle w, w' \rangle = \sum_{\alpha} w_{\alpha} w'_{\alpha}/x_{\alpha}$ and $\|w\|_x^2 = (\sum_{\alpha} w_{\alpha}^2/x_{\alpha})^{1/2}$

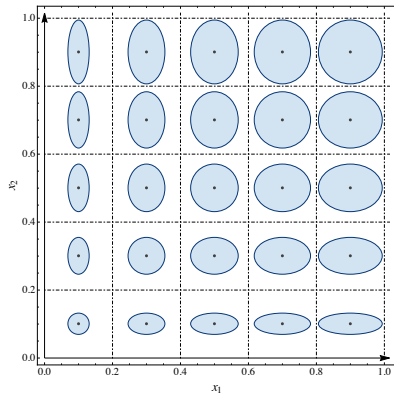


Deformations under different metrics I

Euclidean metric balls



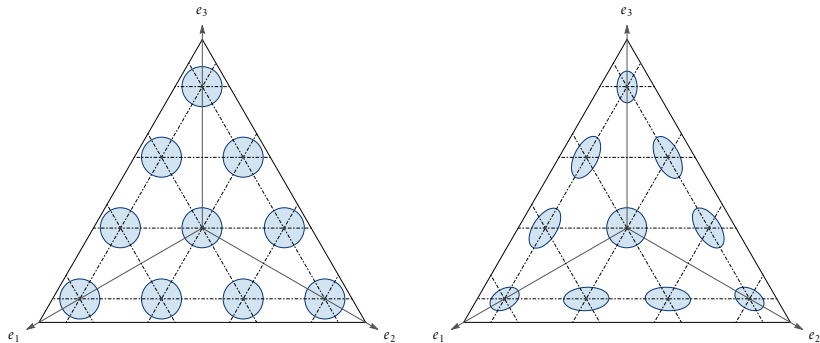
Shahshahani metric balls



Unit balls on the positive orthant of \mathbb{R}^2 under the Euclidean and Shahshahani metrics



Deformations under different metrics II



Unit balls on the unit simplex of \mathbb{R}^2 under the Euclidean and Shahshahani metrics



Duality: payoffs as covectors

The scalar product $\langle w, w' \rangle$ between two vectors $w, w' \in W$ is not to be confused with the action $\langle \omega | w \rangle \equiv \omega(w)$ of a linear functional $\omega \in W^*$ and a vector $w \in W$. In components:

$$\langle \omega | w \rangle = \sum_{\alpha} \omega_{\alpha} w_{\alpha}$$

but this is *not* a scalar product (ω and w live in different spaces)

Terminology: because of duality, linear functionals $\omega \in W^*$ are called **covectors**



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Payoff “vectors” naturally act as linear functionals \implies should be treated as covectors

Examples

1. **Nash equilibrium:** $\sum_{\alpha} v_{\alpha}(x^*)(x_{\alpha} - x_{\alpha}^*) \leq 0$ is a requirement on the action of $v(x^*)$ on the displacement vector $x - x^*$
2. **Potential games:** the derivative $\nabla f = (\partial_1 f, \dots, \partial_n f)$ of a function f naturally acts on vectors to yield the directional derivative $\nabla_z f(x) = \langle \nabla f(x) | z \rangle$
 $\implies v(x) = \nabla f(x)$ behaves as a covector
3. **Positive correlation:** $\langle v(x) | V(x) \rangle \geq 0$ is a requirement on the action of $v(x)$ on the vector of motion $V(x)$



Primal equivalents: payoffs as vectors

Given a Riemannian metric, there is a canonical way to translate covectors into vectors:

Definition

The *primal equivalent* of a covector $\omega \in W^*$ is the unique vector $\omega^\sharp \in W$ such that

$$\langle \omega | w \rangle = \langle \omega^\sharp, w \rangle \quad \text{for all } w \in W.$$

In components, $\omega_\alpha^\sharp = \sum_\beta g_{\alpha\beta}^\sharp(x) \omega_\beta$ where $g^\sharp(x) \equiv g^{-1}(x)$ is the inverse matrix of $g(x)$

Examples

1. In the Euclidean case, $g_{\alpha\beta} = \delta_{\alpha\beta} = g_{\alpha\beta}^\sharp$ so we have

$$\omega_\alpha^\sharp = \omega_\alpha$$

(Source of the standard confusion between scalar products and duality pairings)

2. In the Shahshahani case, $g_{\alpha\beta} = \delta_{\alpha\beta}/x_\beta$, so $g_{\alpha\beta}^\sharp = x_\alpha \delta_{\alpha\beta}$, and we obtain

$$\omega_\alpha^\sharp = x_\alpha \omega_\alpha$$



Extensions to boundary points

The Shahshahani metric explodes at the boundary of \mathcal{X} , so care must be taken there:

Definition

Let \mathcal{K} be a closed convex subset of \mathbb{R}^A with nonempty interior. We say that a Riemannian metric g on \mathcal{K}° is *extendable* to \mathcal{K} if the associated sharp operator g^\sharp admits a (necessarily unique) continuous extension to \mathcal{K} such that

$$T_{\mathcal{K}}(x) \subseteq \text{im } g^\sharp(x) \quad \text{for all } x \in \mathcal{K}.$$

- ▶ The space $\text{im } g^\sharp(x)$ is called the *domain* of g at x and is denoted as $\text{dom } g(x)$.
- ▶ If $\text{dom } g(x) = \mathbb{R}^A$, we say that g is *full-rank extendable*;
instead, if $\text{dom } g(x) = T_{\mathcal{K}}(x)$, we say that g is *minimal-rank extendable*.

Intuition: the extension of g must metrize (at least) all tangent vectors to \mathcal{K} at x .



Extensions to boundary points, cont'd

Why this complicated definition?

Proposition

If g is extendable, there exists a unique scalar product $\langle \cdot, \cdot \rangle_x$ on $\mathbf{dom} g(x)$ such that $\langle w|w' \rangle_{x_k} \rightarrow \langle w|w' \rangle_x$ for all $w, w' \in \mathbf{dom} g(x)$ and for all sequences x_k in \mathcal{K}° with $x_k \rightarrow x$.

In particular, if $\mathcal{K} = \mathbb{R}_+^A$ and g is minimal-rank extendable, we have $g_{\alpha\beta}^\#(x) = 0$ whenever $x_\alpha = 0$ or $x_\beta = 0$.



Extensions to boundary points, cont'd

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Examples

Let \mathcal{K} denote the closed positive orthant of \mathbb{R}^A (= generalized population states). Then:

1. The Euclidean metric $g_{\alpha\beta}(x) = \delta_{\alpha\beta}$ is full-rank extendable (by default).
2. The Shahshahani metric $g_{\alpha\beta}(x) = \delta_{\alpha\beta}/x_\beta$ has $g_{\alpha\beta}^\#(x) = x_\alpha \delta_{\alpha\beta}$, so $\text{dom } g(x) = T_{\mathcal{K}}(x)$

⇒ The Shahshahani metric is minimal-rank extendable and

$$\langle w, w' \rangle_x = \sum_{\alpha \in \text{supp}(x)} w_\alpha w'_\alpha / x_\alpha \quad \text{for all } w, w' \in T_{\mathcal{K}}(x)$$



Tangent projections

Main idea: follow tangent projections of payoffs on \mathcal{X}

Definition

Let g be an extendable Riemannian metric on $\mathcal{K} = \mathbb{R}_+^n$. Then, the *tangent projection* of $w \in \text{dom } g(x)$ to \mathcal{X} at $x \in \mathcal{X}$ is

$$\Pi_x(w) = \arg \min \{ \|w - z\|_x : z \in \text{TC}_{\mathcal{X}}(x) \cap \text{dom } g(x) \}$$



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Remark: minimization is over all tangent vectors $\text{TC}_{\mathcal{X}}(x)$ that are *metrizable*; this need not be the entire tangent cone if g is not full-rank extendable:

Definition

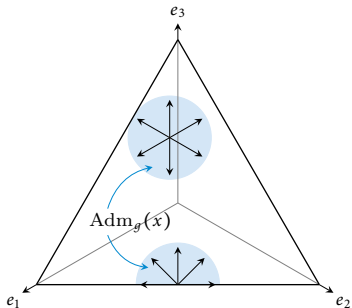
We call this set the *g -admissible tangent cone* of \mathcal{X} at x :

$$\text{Adm}_g(x) = \text{TC}_{\mathcal{X}}(x) \cap \text{dom } g(x)$$

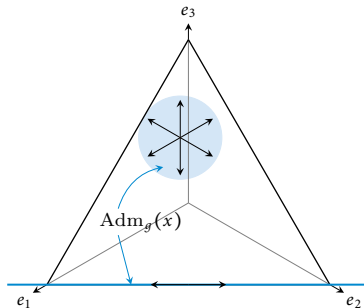


Examples of g -admissible cones

Admissible cones under the Euclidean and Shahshahani metrics.



(e) Euclidean boundary behavior.



(f) Shahshahani boundary behavior.



Coordinate expressions

Case 1: computing tangent projections is straightforward if x is interior:

1. Note that the vector $n(x) = \mathbf{1}^\sharp$ is **normal** to \mathcal{X} at x :

$$\langle n(x), z \rangle = \langle \mathbf{1} | z \rangle = \sum_{\alpha} z_{\alpha} = 0 \quad \text{for all } z \in T_x(x)$$

2. By subtracting the component of w along $n(x)$, we have

$$\Pi_x(w) = w - \text{proj}_{n(x)} w = w - \frac{\langle n(x), w \rangle_x}{\|n(x)\|_x^2} n(x)$$

or, in coordinates:

$$(\Pi_x(w))_{\alpha} = w_{\alpha} - \frac{\sum_{\gamma} w_{\gamma}}{\sum_{\gamma} n_{\gamma}(x)} n_{\alpha}(x)$$

where $n_{\alpha}(x) = \sum_{\beta} g_{\alpha\beta}^{\sharp}(x)$ is the α -th component of $n(x)$.

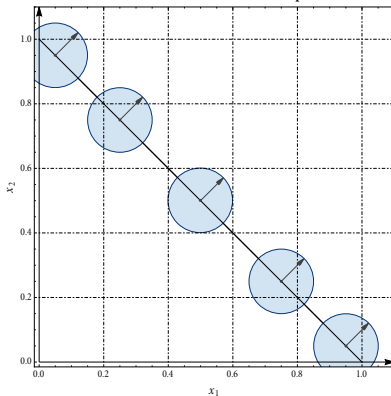
Case 2: construction similar if $\text{dom } g(x) = \mathbb{R}^{\text{supp}(x)}$

Case 3: if $\text{dom } g(x) \not\supseteq \mathbb{R}^{\text{supp}(x)}$, must solve a convex problem.



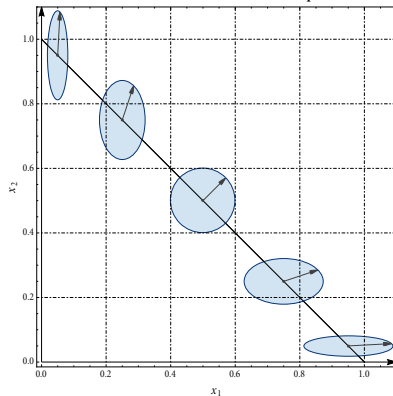
Examples of unit normals

Euclidean structure of the simplex



(g) Euclidean unit normals

Shahshahani structure of the simplex



(h) Shahshahani unit normals



Examples of tangent projections

1. Euclidean projections:

Given $w \in \mathbb{R}^n$, its tangent projection to \mathcal{X} at x is

$$(\Pi_x(w))_\alpha = w_\alpha - |\mathcal{A}(x)|^{-1} \sum_{y \in \mathcal{A}(x)} w_y,$$

where $\mathcal{A}(x)$ is a certain subset of \mathcal{A} that contains $\text{supp}(x)$ – details in Lahkar and Sandholm, 2008.

2. Shahshahani projections:

Given $w \in \mathbb{R}^{\text{supp}(x)}$, its tangent projection to \mathcal{X} at x is

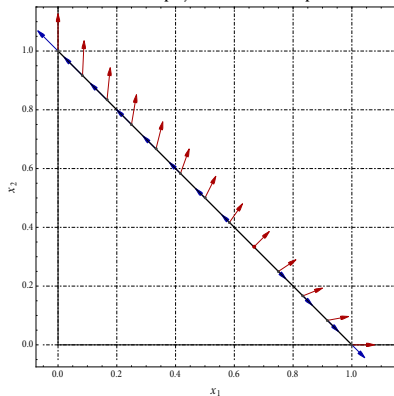
$$(\Pi_x(w))_\alpha = w_\alpha - x_\alpha \sum_y w_y$$

NB: Euclidean and Shahshahani projections do not have the same domain.

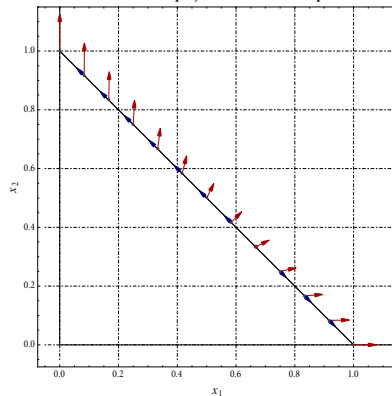


Examples of tangent projections

Euclidean projections on the simplex



Shahshahani projections on the simplex





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Definition and coordinate expressions

Recall the main idea: vectors of motion are aligned to payoffs as closely as possible

This gives the class of *Riemannian game dynamics*

$$\dot{x} = \Pi_x (v^\sharp(x)) \quad (\text{Rie})$$

where $v^\sharp(x)$ is the primal equivalent of the payoff covector $v(x)$ at population state x and $\Pi_x(v^\sharp(x))$ is its tangent projection to \mathcal{X} at x .

In particular, when x is interior or g is minimal-rank extendable, we get (up to a change in time):

$$\dot{x}_\alpha = v_\alpha^\sharp(x) \sum_\gamma n_\gamma(x) - n_\alpha(x) \sum_\gamma v_\gamma^\sharp(x)$$

where $v_\alpha^\sharp(x) = \sum_\beta g_{\alpha\beta}^\sharp(x)$ and $n_\alpha(x) = \sum_\beta g_{\alpha\beta}^\sharp(x)$.



Examples

1. The replicator dynamics:

If g is the Shahshahani metric, $g_{\alpha\beta}^{\sharp}(x) = \delta_{\alpha\beta}x_{\beta}$, so $v_{\alpha}^{\sharp}(x) = x_{\alpha}v_{\alpha}(x)$ and $n_{\alpha}(x) = x_{\alpha}$

\implies

$$\dot{x}_{\alpha} = x_{\alpha} \left[v_{\alpha}(x) - \sum_{\beta} x_{\beta} v_{\beta}(x) \right] \quad (\text{RD})$$

2. The projection dynamics:

If g is the Euclidean metric, $g_{\alpha\beta} = \delta_{\alpha\beta}$, $v_{\alpha}^{\sharp} = v_{\alpha}$ and $n_{\alpha}(x) = 1$. Ultimately, we get:

$$\dot{x}_{\alpha} = \begin{cases} v_{\alpha}(x) - |\mathcal{A}(x)|^{-1} \sum_{\beta \in \mathcal{A}(x)} v_{\beta}(x) & \text{if } \alpha \in \mathcal{A}(x), \\ 0 & \text{otherwise,} \end{cases} \quad (\text{PD})$$

where $\mathcal{A}(x)$ is a subset of \mathcal{A} containing $\text{supp}(x)$ – details in Lahkar and Sandholm (2008).



Examples, cont'd

3. Separable metrics:

Let g be of the form $g_{\alpha\beta}(x) = \delta_{\alpha\beta}/\phi(x_\beta)$ with $\phi: [0, \infty) \rightarrow [0, \infty)$ a *weighting function*. We then get the dynamics:

$$\dot{x}_\alpha = \phi(x_\alpha) \left[v_\alpha(x) \sum_{\beta \in \mathcal{A}} \phi(x_\beta) - \sum_{\beta \in \mathcal{A}} \phi(x_\beta) v_\beta(x) \right]$$

Weighted imitation of success: α -strategists receive revision opportunities with rate $\phi(x_\alpha)$ and switch to β with (conditional) rate $\phi(x_\beta)v_\alpha$



Examples, cont'd

4. Hessian Riemmanian metrics:

Let $h: \mathbb{R}_+^A \rightarrow \mathbb{R}$ be a smooth function s.t. $\text{Hess}(h(x)) > 0$ for all $x > 0$. Then, the *Hessian Riemannian metric* induced by g is

$$g_{\alpha\beta}(x) = \frac{\partial^2 h}{\partial x_\alpha \partial x_\beta}$$

and the induced dynamics take the form:

$$\dot{x}_\alpha = \sum_{\beta \in \text{supp}(x_k)} \left[g_{\alpha\beta}^\sharp(x) - \frac{1}{G^\sharp(x)} g_\alpha^\sharp(x) g_\beta^\sharp(x) \right] v_\beta(x), \quad (\text{HRD})$$

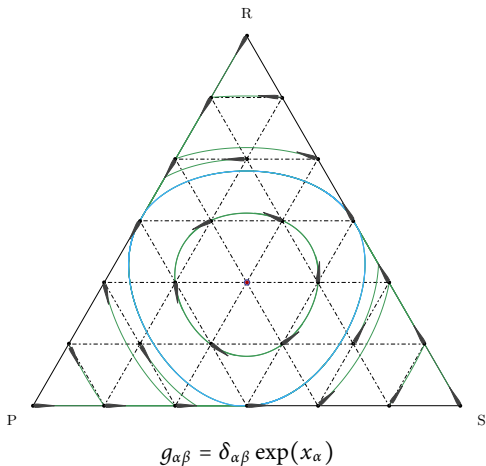
where $g_{\alpha\beta}^\sharp(x) = \text{Hess}(h(x))_{\alpha\beta}^{-1}$, $g_\alpha^\sharp = \sum_\beta g_{\alpha\beta}^\sharp$, and $G^\sharp(x) = \sum_\alpha g_\alpha^\sharp(x)$.

Antecedent: closely related to a class of reinforcement learning dynamics studied by M. and Sandholm (2015).



Phase portraits

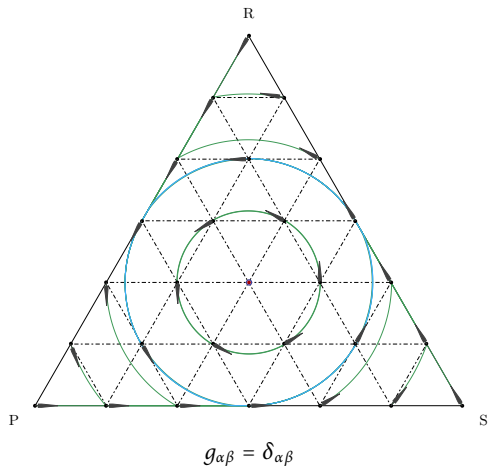
Phase portraits for a Rock-Paper-Scissors game under different metrics:





Phase portraits

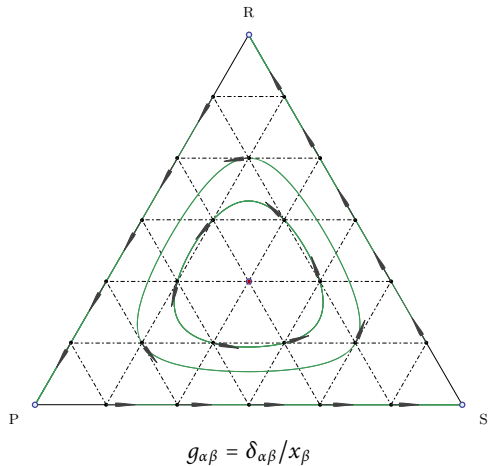
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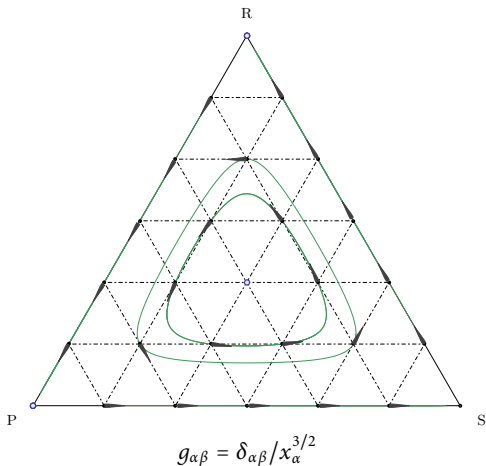
Phase portraits for a Rock-Paper-Scissors game under different metrics:





Phase portraits

Phase portraits for a Rock-Paper-Scissors game under different metrics:





Outline

Background

Geometric Preliminaries

The Dynamics

Analysis



Well-posedness

The first fundamental question is global existence and uniqueness of solutions. This is true if g is Lipschitz continuous:

Proposition

Let g be a Lipschitz extendable Riemannian metric on \mathbb{R}^A . Then, for every $x_0 \in \mathcal{X}$, the induced dynamics admit a unique forward solution $x(t)$ with $x(0) = x_0$.

In particular, if the underlying Riemannian metric is minimal-rank extendable, all solution orbits are smooth and the relative interior of each face of \mathcal{X} is forward invariant: $x_\alpha(t) = 0$ for some $t > 0$ if and only if $x_\alpha(0) = 0$.

Idea of proof.

- ▶ Existence is trivial.
- ▶ Main difficulty for uniqueness is the discontinuity of the motion field at the boundary of \mathcal{X} .
- ▶ The Cojocaru and Jonker theory for projected systems also fails (applies to Hilbert spaces, not Hilbert manifolds).
- ▶ Use ad hoc estimates to show that the system is 1-sided Lipschitz near discontinuity points and apply the differential inclusion results of Kunze (2000). □



Further results

Proposition

Riemannian game dynamics satisfy the positive correlation condition:

$$\langle v(x) | V(x) \rangle \geq 0,$$

with equality if and only if $V(x) = 0$.

Natural: motion field approximates payoff field to the greatest possible extent



Further results

Proposition

Riemannian game dynamics satisfy the positive correlation condition:

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Natural: motion field approximates payoff field to the greatest possible extent

However, the dynamics exhibit an important dichotomy w.r.t. equilibrium play:

Proposition

- ▶ If g is full-rank extendable, $x^* \in \mathcal{X}$ is stationary if and only if it is a Nash equilibrium.
- ▶ If g is minimal-rank extendable, $x^* \in \mathcal{X}$ is stationary if and only if it is a restricted Nash equilibrium (i.e. an equilibrium of a restriction of the game to a subface of \mathcal{X}).

(Explains the difference between the replicator dynamics and the projection dynamics)



Some observations

Some observations before moving further:

- ▶ Dominated strategies may survive under the (Euclidean) projection dynamics (Sandholm et al., 2008).
- ▶ Survival under the projection dynamics is due to discontinuities at the boundary; however, dominated strategies can also survive under *smooth* Riemannian dynamics.

Not all metrics lead to rational outcomes \implies Must refine the class of metrics under study



Distance-like functions

A key element of our analysis is the notion of a distance-like function:

Definition

Let g be an extendable Riemannian metric on \mathcal{K} and let $p \in \mathcal{K}$. A smooth function $D_p: \mathcal{K}^\circ \rightarrow \mathbb{R}$ is called *distance-like* w.r.t. p if:

1. D_p is **positive-definite**: $D_p(x) \geq 0$ for all $x \in \mathcal{K}^\circ$ and $D_p(x_n) \rightarrow 0$ for a sequence $x_k \in \mathcal{K}^\circ$ if and only if $x_k \rightarrow p$.
2. D_p is **radial**:

$$\text{grad} D_p(x) = \sigma(x)(x - p)$$

for some smooth positive function $\sigma: \mathcal{K}^\circ \rightarrow (0, \infty)$.

If g admits distance-like functions, we say that g is *integrable*.

Examples

1. *The Euclidean case*: $D_p(x) = \frac{1}{2} \sum_{\beta} (x_{\beta} - p_{\beta})^2$ is distance-like.
2. *The Shahshahani case*: the Kullback-Leibler divergence $D_p(x) = \sum_{\alpha} p_{\alpha} \log p_{\alpha}/x_{\alpha}$ is distance-like



Elimination of Dominated Strategies

A strategy $p \in \mathcal{X}$ is *dominated* by p' if $\langle v(x) | p - p' \rangle < 0$ for all $x \in \mathcal{X}$.

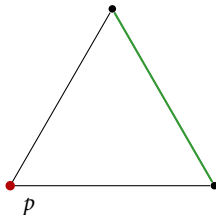


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A strategy $p \in \mathcal{X}$ is *dominated* by p' if $\langle v(x) | p - p' \rangle < 0$ for all $x \in \mathcal{X}$.

A strategy $p \in \mathcal{X}$ becomes *extinct* along $x(t)$ if

$$\min\{x_\alpha(t) : \alpha \in \text{supp}(p)\} \rightarrow 0 \quad \text{as } t \rightarrow \infty$$



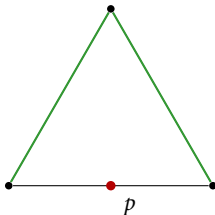


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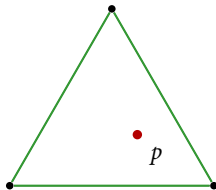


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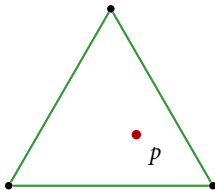


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Theorem

If g is minimal-rank extendable and integrable, dominated strategies become extinct along every interior solution of the induced Riemannian dynamics.



Stability and convergence to equilibrium play

Proposition

Suppose that g is minimal-rank extendable and integrable, and let $x^* \in \mathcal{X}$. If a) $x(t) \rightarrow x^*$ for some interior orbit; or b) x^* is Lyapunov stable, then x^* is a Nash equilibrium.



Stability and convergence to equilibrium play

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Theorem

Suppose that g is minimal-rank extendable and integrable, and let $x^* \in \mathcal{X}$. If x^* is evolutionarily stable, then it is asymptotically stable under the induced Riemannian dynamics; and if it is globally evolutionarily stable, its basin of attraction includes the interior of \mathcal{X} .

(Converse direction: asymptotically stable states are isolated, perfect equilibria)



Permanence

An important evolutionary question is whether all strategies remain above a given threshold in the long run (if not, small fluctuations could lead to extinction of otherwise healthy types):

Definition (Hofbauer and Sigmund, 1998)

A dynamical system on \mathcal{X} is called *permanent* if there exists some $\delta > 0$ such that $\liminf_{t \rightarrow \infty} x_\alpha(t) \geq \delta$ for all $\alpha \in \mathcal{A}$ and for all interior conditions $x(0) \in \mathcal{X}^\circ$.



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Theorem

Suppose that g is minimal-rank extendable and integrable, and assume there is a population state $p \in \mathcal{X}^\circ$ such that

$$\langle v(x^*) | p - x^* \rangle > 0 \quad (\text{P})$$

for all restricted equilibria x^* . Then, the Riemannian game dynamics induced by g are permanent.



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Consequence: If (??) holds, the time-average $\bar{x}(t) = t^{-1} \int_0^t x(s) ds$ of the population's state converges to an interior Nash equilibrium.



Questions?